Optimization problems. CVXPY. Basic linear algebra recap.

### **Daniil Merkulov**

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# **Optimization problems**



# **Example: Transportation problem**

Customer / Source	Arnhem [€/ton]	Gouda [ <b>€</b> /ton]	Demand [tons]
London	n/a	2.5	125
Berlin	2.5	n/a	175
Maastricht	1.6	2.0	225
Amsterdam	1.4	1.0	250
Utrecht	0.8	1.0	225
The Hague	1.4	0.8	200
Supply [tons]	550 tons	700 tons	n/a

minimize: 
$$Cost = \sum_{c \in Customers} \sum_{s \in Sources} T[c, s]x[c, s]$$



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$$\begin{array}{ll} \text{minimize:} & \mathsf{Cost} = \sum_{c \in \mathsf{Customers}} \sum_{s \in \mathsf{Sources}} T[c,s] x[c,s] \\ & \sum_{c \in \mathsf{Customers}} x[c,s] \leq \mathsf{Supply}[s] & \forall s \in \mathsf{Sources} \end{array}$$



# **Example: Transportation problem**

This can be represented in the following graph:

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$$\mathsf{Cost} = \sum_{c \in \mathsf{Customers}} \sum_{s \in \mathsf{Sources}} T[c,s]x[c,s]$$

$$\sum_{c \in \mathsf{Customers}} x[c,s] \le \mathsf{Supply}[s] \qquad \forall s \in \mathsf{Sources}$$

$$\sum_{s \in \mathsf{Sources}} x[c,s] = \mathsf{Demand}[c] \qquad \forall c \in \mathsf{Customers}$$



Figure 1: Graph associated with the problem

## Hardware progress vs Software progress

What would you choose, assuming, that the question posed correctly (you can compile software for any hardware and the problem is the same for both options)? We will consider the time period from 1992 to 2023.

<b>b</b> Hardware	<b>i</b> Software
Solving MIP with an old software on the modern hardware	Solving MIP with a modern software on the old hardware

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It turns out that if you need to solve a MILP, it is better to use an old computer and modern methods than vice versa, the newest computer and methods of the early 1990s!<sup>1</sup>



# **CVXPY** Library



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• User-Friendly Syntax. CVXPY uses a natural mathematical syntax, making it easy to formulate optimization problems.

```
import cvxpy as cp
x = cp.Variable()
objective = cp.Minimize(x**2)
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prob = cp.Problem(objective, constraints)
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- Integration with Scientific Libraries. CVXPY integrates seamlessly with other scientific libraries like NumPy, SciPy, and Pandas, allowing for easy manipulation of data and embedding of optimization problems within larger scientific workflows.

### **CVXPY Example**

```
Minimize the function f(x, y) = x^2 + y^2 s.t. x + y = 1 x - y > 1
import cvxpy as cp
# Define variables
x = cp.Variable()
y = cp.Variable()
# Define objective
objective = cp.Minimize(x**2 + v**2)
# Define constraints
constraints = [x + y == 1, x - y \ge 1]
# Form and solve problem
prob = cp.Problem(objective, constraints)
result = prob.solve()
print(f"Optimal value: {result}")
print(f"Optimal variables: x = {x.value}, y = {y.value}")
```

### **CVXPY** exercises

• CVXPY Examples



## **CVXPY** exercises

- CVXPY Examples
- Exercise



# Basic linear algebra background



#### **Vectors and matrices**

We will treat all vectors as column vectors by default. The space of real vectors of length n is denoted by  $\mathbb{R}^n$ , while the space of real-valued  $m \times n$  matrices is denoted by  $\mathbb{R}^{m \times n}$ . That's it: <sup>2</sup>

$$x = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \quad x^T = \begin{bmatrix} x_1 & x_2 & \dots & x_n \end{bmatrix} \quad x \in \mathbb{R}^n, x_i \in \mathbb{R}$$
(1)

<sup>&</sup>lt;sup>2</sup>A full introduction to applied linear algebra can be found in Introduction to Applied Linear Algebra – Vectors, Matrices, and Least Squares - book by Stephen Boyd & Lieven Vandenberghe, which is indicated in the source. Also, a useful refresher for linear algebra is in Appendix A of the book Numerical Optimization by Jorge Nocedal Stephen J. Wright.

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Similarly, if  $A \in \mathbb{R}^{m \times n}$  we denote transposition as  $A^T \in \mathbb{R}^{n \times m}$ :

$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix} \quad A^{T} = \begin{bmatrix} a_{11} & a_{21} & \dots & a_{m1} \\ a_{12} & a_{22} & \dots & a_{m2} \\ \vdots & \vdots & \ddots & \vdots \\ a_{1n} & a_{2n} & \dots & a_{mn} \end{bmatrix} \quad A \in \mathbb{R}^{m \times n}, a_{ij} \in \mathbb{R}$$

We will write  $x \geq 0$  and  $x \neq 0$  to indicate componentwise relationships

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Figure 2: Equivivalent representations of a vector

A matrix  $A \in \mathbb{S}^n$  is called **positive (negative) definite** if for all  $x \neq 0$ :  $x^T A x > (<)0$ . We denote this as  $A \succ (\prec)0$ . The set of such matrices is denoted as  $\mathbb{S}^n_{++}(\mathbb{S}^n_{--})$ 

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## Matrix product (matmul)

Let A be a matrix of size  $m \times n$ , and B be a matrix of size  $n \times p$ , and let the product AB be:

C = AB

then C is a  $m \times p$  matrix, with element (i, j) given by:

$$c_{ij} = \sum_{k=1}^{n} a_{ik} b_{kj}.$$

This operation in a naive form requires  $O(n^3)$  arithmetical operations, where n is usually assumed as the largest dimension of matrices.

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Is it possible to multiply two matrices faster, than  $\mathcal{O}(n^3)$ ? How about  $\mathcal{O}(n^2)$ ,  $\mathcal{O}(n)$ ?



Let A be a matrix of shape  $m \times n$ , and x be  $n \times 1$  vector, then the *i*-th component of the product:

z = Ax

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$$z_i = \sum_{k=1}^n a_{ik} x_k$$

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 $f \rightarrow \min_{x,y,z}$  Basic linear algebra background

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•  $\langle x, Ay \rangle = \langle A^T x, y \rangle$ 

Norm is a qualitative measure of the smallness of a vector and is typically denoted as ||x||.

The norm should satisfy certain properties:

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3. If  $||x|| = 0$  then  $x = 0$ 

The distance between two vectors is then defined as

$$d(x,y) = \|x - y\|.$$

The most well-known and widely used norm is Euclidean norm:

$$|x||_2 = \sqrt{\sum_{i=1}^n |x_i|^2},$$

which corresponds to the distance in our real life. If the vectors have complex elements, we use their modulus. Euclidean norm, or 2-norm, is a subclass of an important class of p-norms:

$$|x||_{p} = \left(\sum_{i=1}^{n} |x_{i}|^{p}\right)^{1/p}$$



# *p*-norm of a vector

There are two very important special cases. The infinity norm, or Chebyshev norm is defined as the element of the maximal absolute value:

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# $\ensuremath{\textit{p}\text{-norm}}$ of a vector

 $f \rightarrow \min$ 

Basic linear algebra background

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 $L_1$  norm plays a very important role: it all relates to the **compressed sensing** methods that emerged in the mid-00s as one of the most popular research topics. The code for the picture below is available *here:*. Check also *this* video.



Figure 3: Balls in different norms on a plane

### Matrix norms

In some sense there is no big difference between matrices and vectors (you can vectorize the matrix), and here comes the simplest matrix norm **Frobenius** norm:

$$|A||_F = \left(\sum_{i=1}^m \sum_{j=1}^n |a_{ij}|^2\right)^{1/2}$$

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Spectral norm,  $||A||_2$  is one of the most used matrix norms (along with the Frobenius norm).

$$||A||_2 = \sup_{x \neq 0} \frac{||Ax||_2}{||x||_2},$$

It can not be computed directly from the entries using a simple formula, like the Frobenius norm, however, there are efficient algorithms to compute it. It is directly related to the **singular value decomposition** (SVD) of the matrix. It holds

$$||A||_2 = \sigma_1(A) = \sqrt{\lambda_{\max}(A^T A)}$$

where  $\sigma_1(A)$  is the largest singular value of the matrix A.

## Scalar product

The standard scalar (inner) product between vectors x and y from  $\mathbb{R}^n$  is given by

$$\langle x, y \rangle = x^T y = \sum_{i=1}^n x_i y_i = y^T x = \langle y, x \rangle$$

Here  $x_i$  and  $y_i$  are the scalar *i*-th components of corresponding vectors.

#### i Example

Prove, that you can switch the position of a matrix inside a scalar product with transposition:  $\langle x, Ay \rangle = \langle A^T x, y \rangle$  and  $\langle x, yB \rangle = \langle xB^T, y \rangle$ 

#### Matrix scalar product

The standard scalar (inner) product between matrices X and Y from  $\mathbb{R}^{m \times n}$  is given by

$$\langle X, Y \rangle = \operatorname{tr}(X^T Y) = \sum_{i=1}^m \sum_{j=1}^n X_{ij} Y_{ij} = \operatorname{tr}(Y^T X) = \langle Y, X \rangle$$

i Question

Is there any connection between the Frobenious norm  $\|\cdot\|_F$  and scalar product between matrices  $\langle \cdot, \cdot \rangle$ ?



A scalar value  $\lambda$  is an eigenvalue of the n imes n matrix A if there is a nonzero vector q such that

 $Aq = \lambda q.$ 

he vector q is called an eigenvector of A. The matrix A is nonsingular if none of its eigenvalues are zero. The eigenvalues of symmetric matrices are all real numbers, while nonsymmetric matrices may have imaginary eigenvalues. If the matrix is positive definite as well as symmetric, its eigenvalues are all positive real numbers.



# **Eigenvectors and eigenvalues**

#### i Theorem

$$A \succeq (\succ) 0 \Leftrightarrow$$
 all eigenvalues of  $A$  are  $\ge (>) 0$ 

#### i Proof

 $1. \rightarrow$  Suppose some eigenvalue  $\lambda$  is negative and let x denote its corresponding eigenvector. Then

$$Ax = \lambda x \to x^T A x = \lambda x^T x < 0$$

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2.  $\leftarrow$  For any symmetric matrix, we can pick a set of eigenvectors  $v_1, \ldots, v_n$  that form an orthogonal basis of  $\mathbb{R}^n$ . Pick any  $x \in \mathbb{R}^n$ .

$$x^{T}Ax = (\alpha_{1}v_{1} + \ldots + \alpha_{n}v_{n})^{T}A(\alpha_{1}v_{1} + \ldots + \alpha_{n}v_{n})$$
$$= \sum \alpha_{i}^{2}v_{i}^{T}Av_{i} = \sum \alpha_{i}^{2}\lambda_{i}v_{i}^{T}v_{i} \ge 0$$

here we have used the fact that  $v_i^T v_j = 0$ , for  $i \neq j$ .

# **Eigendecomposition (spectral decomposition)**

Suppose  $A \in S_n$ , i.e., A is a real symmetric  $n \times n$  matrix. Then A can be factorized as

 $A = Q\Lambda Q^T,$ 

 $<sup>^{3}</sup>$ A good cheat sheet with matrix decomposition is available at the NLA course website.

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We usually order the eigenvalues as  $\lambda_1 \geq \lambda_2 \geq \ldots \geq \lambda_n$ . We use the notation  $\lambda_i(A)$  to refer to the *i*-th largest eigenvalue of  $A \in S$ . We usually write the largest or maximum eigenvalue as  $\lambda_1(A) = \lambda_{\max}(A)$ , and the least or minimum eigenvalue as  $\lambda_n(A) = \lambda_{\min}(A)$ .

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The largest and smallest eigenvalues satisfy

$$\lambda_{\min}(A) = \inf_{x \neq 0} \frac{x^T A x}{x^T x}, \qquad \lambda_{\max}(A) = \sup_{x \neq 0} \frac{x^T A x}{x^T x}$$



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and consequently  $\forall x \in \mathbb{R}^n$  (Rayleigh quotient):

 $\lambda_{\min}(A)x^Tx \le x^TAx \le \lambda_{\max}(A)x^Tx$ 

The largest and smallest eigenvalues satisfy

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If we use spectral matrix norm, we can get:

$$\kappa(A) = \frac{\sigma_{\max}(A)}{\sigma_{\min}(A)}$$

If, moreover, 
$$A \in \mathbb{S}^n_{++}$$
:  $\kappa(A) = \frac{\lambda_{\max}(A)}{\lambda_{\min}(A)}$ 

 $f \rightarrow \min_{x,y,z}$  Basic linear algebra background

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This factorization is called the **singular value decomposition (SVD)** of A. The columns of U are called left singular vectors of A, the columns of V are right singular vectors, and the numbers  $\sigma_i$  are the singular values. The singular value decomposition can be written as

$$A = \sum_{i=1}^{r} \sigma_i u_i v_i^T,$$

where  $u_i \in \mathbb{R}^m$  are the left singular vectors, and  $v_i \in \mathbb{R}^n$  are the right singular vectors.

#### 1 Question

Suppose, matrix  $A \in \mathbb{S}^n_{++}$ . What can we say about the connection between its eigenvalues and singular values?



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#### i Question

How do the singular values of a matrix relate to its eigenvalues, especially for a symmetric matrix?



Simple, yet very interesting decomposition is Skeleton decomposition, which can be written in two forms:

$$A = UV^T \quad A = \hat{C}\hat{A}^{-1}\hat{R}$$

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• Model reduction, data compression, and speedup of computations in numerical analysis: given rank-r matrix with  $r \ll n, m$  one needs to store  $\mathcal{O}((n+m)r) \ll nm$  elements.





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Figure 4: Illustration of Skeleton decomposition

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- Feature extraction in machine learning, where it is also known as matrix factorization
- All applications where SVD applies, since Skeleton decomposition can be transformed into truncated SVD form.



Figure 4: Illustration of Skeleton decomposition

# **Canonical tensor decomposition**

One can consider the generalization of Skeleton decomposition to the higher order data structure, like tensors, which implies representing the tensor as a sum of r primitive tensors.



Figure 5: Illustration of Canonical Polyadic decomposition

#### i Example

Note, that there are many tensor decompositions: Canonical, Tucker, Tensor Train (TT), Tensor Ring (TR), and others. In the tensor case, we do not have a straightforward definition of *rank* for all types of decompositions. For example, for TT decomposition rank is not a scalar, but a vector.

The determinant and trace can be expressed in terms of the eigenvalues

$$\det A = \prod_{i=1}^n \lambda_i, \qquad \mathrm{tr} A = \sum_{i=1}^n \lambda_i$$

The determinant has several appealing (and revealing) properties. For instance,

• detA = 0 if and only if A is singular;

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Don't forget about the cyclic property of a trace for arbitrary matrices A, B, C, D (assuming, that all dimensions are consistent):

$$tr(ABCD) = tr(DABC) = tr(CDAB) = tr(BCDA)$$



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$$tr(ABCD) = tr(DABC) = tr(CDAB) = tr(BCDA)$$

#### i Question

How does the determinant of a matrix relate to its invertibility?
The first-order Taylor approximation, also known as the linear approximation, is centered around some point  $x_0$ . If  $f : \mathbb{R}^n \to \mathbb{R}$  is a differentiable function, then its first-order Taylor approximation is given by:

$$f_{x_0}^{I}(x) = f(x_0) + \nabla f(x_0)^{T}(x - x_0)$$

Where:

•  $f(x_0)$  is the value of the function at the point  $x_0$ .

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It is very usual to replace the f(x) with  $f_{x_0}^I(x)$  near the point  $x_0$  for simple analysis of some approaches.



Figure 6: First order Taylor approximation near the point  $x_0$ 

## Second-order Taylor approximation

The second-order Taylor approximation, also known as the quadratic approximation, includes the curvature of the function. For a twice-differentiable function  $f : \mathbb{R}^n \to \mathbb{R}$ , its second-order Taylor approximation centered at some point  $x_0$  is:

$$f_{x_0}^{II}(x) = f(x_0) + \nabla f(x_0)^T (x - x_0) + \frac{1}{2} (x - x_0)^T \nabla^2 f(x_0) (x - x_0)$$

Where  $\nabla^2 f(x_0)$  is the Hessian matrix of f at the point  $x_0$ .



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Where  $\nabla^2 f(x_0)$  is the Hessian matrix of f at the point  $x_0$ . When using the linear approximation of the function is not sufficient one can consider replacing the f(x) with  $f_{x_0}^{II}(x)$  near the point  $x_0$ . In general, Taylor approximations give us a way to locally approximate functions. The first-order approximation is a plane tangent to the function at the point  $x_0$ , while the second-order approximation includes the curvature and is represented by a parabola. These approximations are especially useful in optimization and numerical methods because they provide a tractable way to work with complex functions.



Figure 7: Second order Taylor approximation near the point  $x_0$ 

• Linear Least Squares



- Linear Least Squares
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- Problems



- Linear Least Squares
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- How to calculate minimum and maximum eigenvalue of the hessian matrix of linear least squares problem? What about binary logistic regression?